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# Second-order smoothing approximation to $l_1$ exact penalty function for nonlinear constrained optimization problems

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## Abstract

In this paper, a second-order smoothing approximation to the  $l_1$  exact penalty function for nonlinear constrained optimization problems is presented. Error estimations are obtained among the optimal objective function values of the smoothed penalty problem, of the nonsmooth penalty problem and of the original optimization problem. Based on the smoothed penalty problem, an algorithm that has better convergence is presented. Numerical examples illustrate that this algorithm is efficient in solving nonlinear constrained optimization problems.

**Mathematics Subject Classification:** 90C30; 65K05

**Keywords:** Exact penalty function; Smoothing method; Constrained optimization; Approximate optimal solution

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# 1 Introduction

Consider the nonlinear inequality constrained optimization problem:

$$(P) \quad \min f(x) \\ \text{s.t. } g_i(x) \leq 0, \quad i = 1, 2, \dots, m,$$

where  $f, g_i : R^n \rightarrow R$ ,  $i \in I = \{1, 2, \dots, m\}$  are twice continuously differentiable functions and  $X_0 = \{x \in R^n \mid g_i(x) \leq 0, \quad i = 1, 2, \dots, m\}$  is the feasible set to (P).

To solve (P), many exact penalty function methods have been introduced in literatures, see, [1, 3, 4, 5, 7, 13, 24]. In 1967, Zangwill [24] first the classical  $l_1$  exact penalty function as follows:

$$F(x, \rho) = f(x) + \rho \sum_{i=1}^m \max\{g_i(x), 0\}, \quad (1)$$

where  $\rho > 0$  is a penalty parameter.

Note that penalty function (1) is not a smooth function. The obvious difficulty with the exact penalty function is that it is non-differentiable, which prevents the use of efficient minimization algorithms and may cause some numerical instability problems in its implementation. In order to avoid the drawback related to the nondifferentiability, the smoothing methods of the exact penalty functions attracts much attention, see, [2, 8, 9, 10, 11, 12, 17, 18, 19, 20, 21, 23, 25]. Chen et al. [2] introduced a smooth function to approximate the classical  $l_1$  penalty function by integrating the sigmoid function  $1/(1 + e^{-\alpha t})$ . Lian [8] and Wu et al. [18] proposed a smoothing approximation to  $l_1$  exact penalty function for inequality constrained optimization. Pinar et al. [12] also proposed a smoothing approximation to  $l_1$  exact penalty function and an  $\epsilon$ -optimal minimum can be obtained by solving the smoothed penalty problem. Xu et al. [20] discussed a second-order differentiability smoothing to the classical  $l_1$  exact penalty function for constrained optimization problems. Meng et al. [10] introduced a smoothing of the square-root exact penalty function for inequality constrained optimization. However, little attention has been paid to smoothing the exact penalty function in terms of second-order differentiability. So, here we present a second-order smoothing approximation to the  $l_1$  exact penalty function (1), and based on the smoothed penalty function ob-

tained thereafter an algorithm for solving nonlinear constrained optimization problems is given in this paper.

The rest of this paper is organized as follows. In Section 2, we introduce a smoothing function for the classical  $l_1$  exact penalty function and some fundamental properties of the smoothing function. In Section 2, the algorithm based on the smoothed penalty function is proposed and its global convergence is presented, with some numerical examples given. Finally, conclusions are given in Section 3.

## 2 Second-order smoothing penalty function

Let  $q(t) = \max\{t, 0\}$ . Then, the penalty function (1) is turned into

$$G(x, \rho) = f(x) + \rho \sum_{i=1}^m q(g_i(x)), \quad (2)$$

where  $\rho > 0$ . The corresponding penalty optimization problem to  $G(x, \rho)$  is defined as

$$(P_\rho) \quad \min G(x, \rho), \quad \text{s.t. } x \in R^n.$$

In order to  $q(t)$ , we define function  $q_\epsilon(t) : R^1 \rightarrow R^1$  as

$$q_\epsilon(t) = \begin{cases} 0 & \text{if } t < 0, \\ \frac{t^3}{9\epsilon^2} & \text{if } 0 \leq t < \epsilon, \\ t + \frac{2}{3}\epsilon e^{-\frac{t}{\epsilon}+1} - \frac{14\epsilon}{9} & \text{if } t \geq \epsilon, \end{cases}$$

where  $\epsilon > 0$  is a smoothing parameter.

**Remark 1.** Obviously,  $q_\epsilon(t)$  has the following attractive properties:

(i) For any  $\epsilon > 0$ ,  $q_\epsilon(t)$  is twice continuously differentiable on  $R$ , where

$$q'_\epsilon(t) = \begin{cases} 0 & \text{if } t < 0, \\ \frac{t^2}{3\epsilon^2} & \text{if } 0 \leq t < \epsilon, \\ 1 - \frac{2}{3}e^{-\frac{t}{\epsilon}+1} & \text{if } t \geq \epsilon, \end{cases}$$

and

$$q''_\epsilon(t) = \begin{cases} 0 & \text{if } t < 0, \\ \frac{2t}{3\epsilon^2} & \text{if } 0 \leq t < \epsilon, \\ \frac{2}{3\epsilon} e^{-\frac{t}{\epsilon}+1} & \text{if } t \geq \epsilon. \end{cases}$$

(ii)  $\lim_{\epsilon \rightarrow 0} q_\epsilon(t) = q(t)$ .

(iii)  $q_\epsilon(t)$  is convex and monotonically increasing in  $t$  for any given  $\epsilon > 0$ .

Property (iii) follow from (i) immediately.

Suppose that  $f$  and  $g_i (i = 1, 2, \dots, m)$  are second-order continuously differentiable. Consider the penalty function for (P) given by

$$G_\epsilon(x, \rho) = f(x) + \rho \sum_{i=1}^m q_\epsilon(g_i(x)). \quad (3)$$

Clearly,  $G_\epsilon(x, \rho)$  is second-order continuously differentiable on  $R^n$ . Applying (3), the following penalty problem for (P) is obtained

$$(SP_{\rho, \epsilon}) \quad \min G_\epsilon(x, \rho), \quad \text{s.t. } x \in R^n.$$

Now, the relationship between  $(P_\rho)$  and  $(SP_{\rho, \epsilon})$  is studied.

**Lemma 2.1.** *For any given  $x \in R^n$ ,  $\epsilon > 0$  and  $\rho > 0$ , we have*

$$0 \leq G(x, \rho) - G_\epsilon(x, \rho) \leq \frac{14m\rho\epsilon}{9}. \quad (4)$$

*Proof.* For  $x \in R^n$  and  $i \in I$ , by the definition of  $q_\epsilon(t)$ , we have

$$q(g_i(x)) - q_\epsilon(g_i(x)) = \begin{cases} 0 & \text{if } g_i(x) < 0, \\ g_i(x) - \frac{g_i(x)^3}{9\epsilon^2} & \text{if } 0 \leq g_i(x) < \epsilon, \\ \frac{14\epsilon}{9} - \frac{2}{3}\epsilon e^{-\frac{g_i(x)}{\epsilon}+1} & \text{if } g_i(x) \geq \epsilon. \end{cases}$$

That is,

$$0 \leq q(g_i(x)) - q_\epsilon(g_i(x)) \leq \frac{14\epsilon}{9}, \quad i = 1, 2, \dots, m.$$

Thus,

$$0 \leq \sum_{i=1}^m q(g_i(x)) - \sum_{i=1}^m q_\epsilon(g_i(x)) \leq \frac{14m\epsilon}{9},$$

which implies

$$0 \leq \rho \sum_{i=1}^m q(g_i(x)) - \rho \sum_{i=1}^m q_\epsilon(g_i(x)) \leq \frac{14m\rho\epsilon}{9}.$$

Therefore,

$$0 \leq \left\{ f(x) + \rho \sum_{i=1}^m q(g_i(x)) \right\} - \left\{ f(x) + \rho \sum_{i=1}^m q_\epsilon(g_i(x)) \right\} \leq \frac{14m\rho\epsilon}{9},$$

that is,

$$0 \leq G(x, \rho) - G_\epsilon(x, \rho) \leq \frac{14m\rho\epsilon}{9}.$$

The proof completes.  $\square$

A direct result of Lemma 2.1 is given as follows.

**Corollary 2.2.** *Let  $\{\epsilon_j\} \rightarrow 0$  be a sequence of positive numbers and assume  $x^j$  is a solution to  $(SP_{\rho,\epsilon})$  for some given  $\rho > 0$ . Let  $x'$  be an accumulation point of the sequence  $\{x^j\}$ . Then  $x'$  is an optimal solution to  $(P_\rho)$ .*

**Definition 2.3.** *For  $\epsilon > 0$ , a point  $x_\epsilon \in R^n$  is called  $\epsilon$ -feasible solution to  $(P)$  if  $g_i(x_\epsilon) \leq \epsilon$ ,  $\forall i \in I$ .*

**Definition 2.4.** *For  $\epsilon > 0$ , a point  $x_\epsilon \in X_0$  is called  $\epsilon$ -approximate optimal solution to  $(P)$  if*

$$|f^* - f(x_\epsilon)| \leq \epsilon,$$

where  $f^*$  is the optimal objective value of  $(P)$ .

**Theorem 2.5.** *Let  $x^*$  be an optimal solution of problem  $(P_\rho)$  and  $x'$  be an optimal solution to  $(SP_{\rho,\epsilon})$  for some  $\rho > 0$  and  $\epsilon > 0$ . Then,*

$$0 \leq G(x^*, \rho) - G_\epsilon(x', \rho) \leq \frac{14m\rho\epsilon}{9}. \quad (5)$$

*Proof.* From Lemma 2.1, for  $\rho > 0$ , we have that

$$\begin{aligned} 0 &\leq G(x^*, \rho) - G_\epsilon(x^*, \rho) \leq \frac{14m\rho\epsilon}{9}, \\ 0 &\leq G(x', \rho) - G_\epsilon(x', \rho) \leq \frac{14m\rho\epsilon}{9}. \end{aligned}$$

Under the assumption that  $x^*$  is an optimal solution to  $(P_\rho)$  and  $x'$  is an optimal solution to  $(SP_{\rho,\epsilon})$ , we get

$$\begin{aligned} G(x^*, \rho) &\leq G(x', \rho), \\ G_\epsilon(x', \rho) &\leq G_\epsilon(x^*, \rho). \end{aligned}$$

Therefore, we obtain that

$$\begin{aligned} 0 &\leq G(x^*, \rho) - G_\epsilon(x^*, \rho) \leq G(x^*, \rho) - G_\epsilon(x', \rho) \\ &\leq G(x', \rho) - G_\epsilon(x', \rho) \leq \frac{14m\rho\epsilon}{9}. \end{aligned}$$

That is,

$$0 \leq G(x^*, \rho) - G_\epsilon(x', \rho) \leq \frac{14m\rho\epsilon}{9}.$$

This completes the proof.  $\square$

**Lemma 2.6 ([19]).** *Suppose that  $x^*$  is an optimal solution to  $(P_\rho)$ . If  $x^*$  is feasible to  $(P)$ , then it is an optimal solution to  $(P)$ .*

**Theorem 2.7.** *Suppose that  $x^*$  satisfies the conditions in Lemma 2.6 and  $x'$  be an optimal solution to  $(SP_{\rho,\epsilon})$  for some  $\rho > 0$  and  $\epsilon > 0$ . If  $x'$  is  $\epsilon$ -feasible to  $(P)$ . Then,*

$$0 \leq f(x^*) - f(x') \leq \frac{5m\rho\epsilon}{3}, \quad (6)$$

that is,  $x'$  is an approximate optimal solution to  $(P)$ .

*Proof.* Since  $x'$  is  $\epsilon$ -feasible to  $(P)$ , it follows that

$$\sum_{i=1}^m q_\epsilon(g_i(x')) \leq \frac{m\epsilon}{9}.$$

As  $x^*$  is a feasible solution to  $(P)$ , we have

$$\sum_{i=1}^m q(g_i(x^*)) = 0.$$

By Theorem 2.5, we get

$$0 \leq \left\{ f(x^*) + \rho \sum_{i=1}^m q(g_i(x^*)) \right\} - \left\{ f(x') + \rho \sum_{i=1}^m q_\epsilon(g_i(x')) \right\} \leq \frac{14m\rho\epsilon}{9}.$$

Thus,

$$\rho \sum_{i=1}^m q_{\epsilon}(g_i(x')) \leq f(x^*) - f(x') \leq \rho \sum_{i=1}^m q_{\epsilon}(g_i(x')) + \frac{14m\rho\epsilon}{9}.$$

That is,

$$0 \leq f(x^*) - f(x') \leq \frac{5m\rho\epsilon}{3}.$$

By Lemma 2.6,  $x^*$  is actually an optimal solution to (P). Thus  $x'$  is an approximate optimal solution to (P). This completes the proof.  $\square$

Theorem 2.5 show that an approximate solution to  $(SP_{\rho,\epsilon})$  is also an approximate solution to  $(P_{\rho})$  when the error  $\epsilon$  is sufficiently small. By Theorem 2.7, an optimal solution to  $(SP_{\rho,\epsilon})$  is an approximate optimal solution to (P) if it is  $\epsilon$ -feasible to (P).

**Definition 2.8.** For  $x^* \in R^n$ , we call  $y^* \in R^m$  a Lagrange multiplier vector corresponding to  $x^*$  if and only if  $x^*$  and  $y^*$  satisfy that

$$\nabla f(x^*) + \sum_{i=1}^m y_i^* \nabla g_i(x^*) = 0, \quad (7)$$

$$y_i^* g_i(x^*) = 0, \quad g_i(x^*) \leq 0, \quad y_i^* \geq 0, \quad i \in I. \quad (8)$$

**Theorem 2.9.** Let  $f$  and  $g_i$ ,  $i \in I$  in (P) are convex. Let  $x^*$  be an optimal solution of (P) and  $y^* \in R^m$  a Lagrange multiplier vector corresponding to  $x^*$ . Then for some  $\epsilon > 0$ ,

$$G(x^*, \rho) - G_{\epsilon}(x, \rho) \leq \frac{14m\rho\epsilon}{9}, \quad \forall x \in R^n, \quad (9)$$

provided that  $\rho \geq y_i^*$ ,  $i = 1, 2, \dots, m$ .

*Proof.* By the convexity of  $f$  and  $g_i$ ,  $i = 1, 2, \dots, m$ , we have

$$\begin{aligned} f(x) &\geq f(x^*) + \nabla f(x^*)^T(x - x^*), \quad \forall x \in R^n, \\ g_i(x) &\geq g_i(x^*) + \nabla g_i(x^*)^T(x - x^*), \quad \forall x \in R^n. \end{aligned} \quad (10)$$

Since  $x^*$  is an optimal solution of (P) and  $y^*$  is a Lagrange multiplier vector corresponding to  $x^*$ , by (7), (8), (9) and (10), we have

$$\begin{aligned}
f(x) &\geq f(x^*) + \nabla f(x^*)^T(x - x^*) \\
&= f(x^*) - \sum_{i=1}^m y_i^* \nabla g_i(x^*)^T(x - x^*) \\
&\geq f(x^*) - \sum_{i=1}^m y_i^* (g_i(x) - g_i(x^*)) \\
&= f(x^*) - \sum_{i=1}^m y_i^* g_i(x).
\end{aligned}$$

Since  $g_i(x) \leq g_i^+(x)$  ( $g_i^+(x) = \max\{0, g_i(x)\}$ ,  $i \in I$ ), we have

$$\begin{aligned}
G(x, \rho) = f(x) + \rho \sum_{i=1}^m g_i^+(x) &\geq f(x^*) - \sum_{i=1}^m y_i^* g_i(x) + \rho \sum_{i=1}^m g_i^+(x) \\
&\geq f(x^*) + \sum_{i=1}^m (\rho - y_i^*) g_i^+(x).
\end{aligned}$$

Thus, for  $\rho \geq y_i^*$ ,  $i = 1, 2, \dots, m$ , we get  $G(x, \rho) \geq f(x^*)$ . Since  $x^*$  is feasible, then  $f(x^*) = G(x^*, \rho)$  and by Lemma 2.1, we have

$$\begin{aligned}
G(x^*, \rho) - G_\epsilon(x, \rho) &= G(x^*, \rho) - G(x, \rho) + G(x, \rho) - G_\epsilon(x, \rho) \\
&= f(x^*) - G(x, \rho) + G(x, \rho) - G_\epsilon(x, \rho) \\
&\leq G(x, \rho) - G_\epsilon(x, \rho) \leq \frac{14m\rho\epsilon}{9}.
\end{aligned}$$

This completes the proof.  $\square$

By Theorem 2.9, when the parameter  $\rho$  is sufficiently large, an approximate optimal solution to  $(SP_{\rho, \epsilon})$  is an approximate optimal solution to (P), where (P) is a convex problem. Therefore, we may obtain an approximate optimal solution to (P) by computing an approximate optimal solution to  $(SP_{\rho, \epsilon})$ .  
sectionAlgorithm and numerical examples

In this section, using the smoothed penalty function  $G_\epsilon(x, \rho)$ , we propose an algorithm to solve nonlinear constrained optimization problems, defined as Algorithm I



**Algorithm I**

Step 1: Choose  $x^0$ ,  $\epsilon > 0$ ,  $\epsilon_0 > 0$ ,  $\rho_0 > 0$ ,  $0 < \eta < 1$  and  $N > 1$ , let  $j = 0$  and go to Step 2.

Step 2: Use  $x^j$  as the starting point to solve

$$(SP_{\rho_j, \epsilon_j}) \quad \min_{x \in R^n} G_{\epsilon_j}(x, \rho_j) = f(x) + \rho_j \sum_{i=1}^m q_{\epsilon_j}(g_i(x)).$$

Let  $x^{j+1}$  be the optimal solution obtained ( $x^{j+1}$  is obtained by a quasi-Newton method).

Step 3: If  $x^{j+1}$  is  $\epsilon$ -feasible to (P), then stop and we have obtained an approximate solution  $x^{j+1}$  of (P). Otherwise, let  $\rho_{j+1} = N\rho_j$ ,  $\epsilon_{j+1} = \eta\epsilon_j$  and  $j = j + 1$ , then go to Step 2.

**Remark 2.** In this Algorithm I, as  $N > 1$  and  $0 < \eta < 1$ , the sequence  $\{\epsilon_j\} \rightarrow 0$  ( $j \rightarrow +\infty$ ) and the sequence  $\{\rho_j\} \rightarrow +\infty$  ( $j \rightarrow +\infty$ ).

In practice, it is difficult to compute  $x^{j+1} \in \arg \min_{x \in R^n} G_{\epsilon_j}(x, \rho_j)$ . We generally look for the local minimizer or stationary point of  $G_{\epsilon_j}(x, \rho_j)$  by computing  $x^{j+1}$  such that  $\nabla G_{\epsilon_j}(x, \rho_j) = 0$ .

For  $x \in R^n$ , we define

$$\begin{aligned} I^0(x) &= \{i \mid g_i(x) < 0, \quad i \in I\}, \\ I_\epsilon^+(x) &= \{i \mid g_i(x) \geq \epsilon, \quad i \in I\}, \\ I_\epsilon^-(x) &= \{i \mid 0 \leq g_i(x) < \epsilon, \quad i \in I\}. \end{aligned}$$

Then, the following result is obtained.

**Theorem 2.10.** *Assume that  $\lim_{\|x\| \rightarrow +\infty} f(x) = +\infty$ . Let  $\{x^j\}$  be the sequence generated by Algorithm I. Suppose that the sequence  $\{G_{\epsilon_j}(x^j, \rho_j)\}$  is bounded. Then  $\{x^j\}$  is bounded and any limit point  $x^*$  of  $\{x^j\}$  is feasible to (P), and satisfies*

$$\lambda \nabla f(x^*) + \sum_{i \in I} \mu_i \nabla g_i(x^*) = 0, \quad (11)$$

where  $\lambda \geq 0$  and  $\mu_i \geq 0$ ,  $i = 1, 2, \dots, m$ .

*Proof.* First, we will prove that  $\{x^j\}$  is bounded. Note that

$$G_{\epsilon_j}(x^j, \rho_j) = f(x^j) + \rho_j \sum_{i=1}^m q_{\epsilon_j}(g_i(x^j)), \quad j = 0, 1, 2, \dots, \quad (12)$$

and by the definition of  $q_\epsilon(t)$ , we have

$$\sum_{i=1}^m q_{\epsilon_j} (g_i(x^j)) \geq 0. \quad (13)$$

Suppose to the contrary that  $\{x^j\}$  is unbounded. Without loss of generality, we assume that  $\|x^j\| \rightarrow +\infty$  as  $j \rightarrow +\infty$ . Then,  $\lim_{j \rightarrow +\infty} f(x^j) = +\infty$  and from (12) and (13), we have

$$G_{\epsilon_j}(x^j, \rho_j) \geq f(x^j) \rightarrow +\infty, \quad \rho_j > 0, \quad j = 0, 1, 2, \dots,$$

which results in a contradiction since the sequence  $\{G_{\epsilon_j}(x^j, \rho_j)\}$  is bounded. Thus  $\{x^j\}$  is bounded.

We show next that any limit point  $x^*$  of  $\{x^j\}$  is feasible to (P). Without loss of generality, we assume that  $\lim_{j \rightarrow +\infty} x^j = x^*$ . Suppose that  $x^*$  is not feasible to (P). Then there exists some  $i \in I$  such that  $g_i(x^*) \geq \alpha > 0$ . Note that

$$\begin{aligned} G_{\epsilon_j}(x^j, \rho_j) &= f(x^j) + \rho_j \sum_{i \in I_{\epsilon_j}^+(x^j)} \left( g_i(x^j) + \frac{2}{3} \epsilon_j e^{-\frac{g_i(x^j)}{\epsilon_j} + 1} - \frac{14\epsilon_j}{9} \right) \\ &\quad + \rho_j \sum_{i \in I_{\epsilon_j}^-(x^j)} \frac{g_i(x^j)^3}{9\epsilon_j^2}. \end{aligned} \quad (14)$$

If  $j \rightarrow +\infty$ , then for any sufficiently large  $j$ , the set  $\{i \mid g_i(x^j) \geq \alpha\}$  is not empty. Because  $I$  is finite, then there exists an  $i_0 \in I$  that satisfies  $g_{i_0}(x^j) \geq \alpha$ . If  $j \rightarrow +\infty$ ,  $\rho_j \rightarrow +\infty$ ,  $\epsilon_j \rightarrow 0$ , it follows from (14) that  $G_{\epsilon_j}(x^j, \rho_j) \rightarrow +\infty$ , which contradicts the assumption that  $\{G_{\epsilon_j}(x^j, \rho_j)\}$  is bounded. Therefore,  $x^*$  is feasible to (P).

Finally, we show that (11) holds. By Step 2 in Algorithm I,  $\nabla G_{\epsilon_j}(x^j, \rho_j) = 0$ , that is

$$\begin{aligned} \nabla f(x^j) + \rho_j \sum_{i \in I_{\epsilon_j}^+(x^j)} \left( 1 - \frac{2}{3} e^{-\frac{g_i(x^j)}{\epsilon_j} + 1} \right) \nabla g_i(x^j) \\ + \rho_j \sum_{i \in I_{\epsilon_j}^-(x^j)} \frac{1}{3\epsilon_j^2} g_i(x^j)^2 \nabla g_i(x^j) = 0. \end{aligned} \quad (15)$$

For  $j = 1, 2, \dots$ , let

$$\gamma_j = 1 + \sum_{i \in I_{\epsilon_j}^+(x^j)} \rho_j \left( 1 - \frac{2}{3} e^{-\frac{g_i(x^j)}{\epsilon_j} + 1} \right) + \sum_{i \in I_{\epsilon_j}^-(x^j)} \frac{\rho_j}{3\epsilon_j^2} g_i(x^j)^2. \quad (16)$$

Then  $\gamma_j > 1$ . From (15), we have

$$\begin{aligned} \frac{1}{\gamma_j} \nabla f(x^j) + \sum_{i \in I_{\epsilon_j}^+(x^j)} \frac{\rho_j \left(1 - \frac{2}{3} e^{-\frac{g_i(x^j)}{\epsilon_j} + 1}\right)}{\gamma_j} \nabla g_i(x^j) \\ + \sum_{i \in I_{\epsilon_j}^-(x^j)} \frac{\rho_j \epsilon_j^{-2}}{3\gamma_j} g_i(x^j)^2 \nabla g_i(x^j) = 0. \end{aligned} \quad (17)$$

Let

$$\begin{aligned} \lambda^j &= \frac{1}{\gamma_j}, \\ \mu_i^j &= \frac{\rho_j \left(1 - \frac{2}{3} e^{-\frac{g_i(x^j)}{\epsilon_j} + 1}\right)}{\gamma_j}, \quad i \in I_{\epsilon_j}^+(x^j), \\ \mu_i^j &= \frac{\rho_j \epsilon_j^{-2}}{3\gamma_j} g_i(x^j)^2, \quad i \in I_{\epsilon_j}^-(x^j), \\ \mu_i^j &= 0, \quad i \in I \setminus \left(I_{\epsilon_j}^+(x^j) \cup I_{\epsilon_j}^-(x^j)\right). \end{aligned}$$

Then we have

$$\begin{aligned} \lambda^j + \sum_{i \in I} \mu_i^j &= 1, \quad \forall j, \\ \mu_i^j &\geq 0, \quad i \in I, \quad \forall j. \end{aligned} \quad (18)$$

When  $j \rightarrow \infty$ , we have that  $\lambda^j \rightarrow \lambda \geq 0$ ,  $\mu_i^j \rightarrow \mu_i \geq 0$ ,  $\forall i \in I$ . By (17) and (18), as  $j \rightarrow +\infty$ , we have

$$\begin{aligned} \lambda \nabla f(x^*) + \sum_{i \in I} \mu_i \nabla g_i(x^*) &= 0, \\ \lambda + \sum_{i \in I} \mu_i &= 1. \end{aligned}$$

For  $i \in I^0(x^*)$ , as  $j \rightarrow +\infty$ , we get  $\mu_i^j \rightarrow 0$ . Therefore,  $\mu_i = 0$ ,  $\forall i \in I^0(x^*)$ . So, (11) holds, and this completes the proof.  $\square$

Now, we will solve some nonlinear constrained optimization problems with Algorithm I on MATLAB. In each of the following examples, the MATLAB 7.12 subroutine *fmincon* is used to obtain the local minima of problem  $(SP_{\rho_j, \epsilon_j})$ .

The numerical results of each example are presented in the following tables. It is shown that Algorithm I yield some approximate solutions that have a better objective function value in comparison with some other algorithms.

**Example 1.** Consider the example in [8],

$$\begin{aligned}
 \text{(P3.1)} \quad \min \quad & f(x) = x_1^2 + x_2^2 + 2x_3^2 + x_4^2 - 5x_1 - 5x_2 - 21x_3 + 7x_4 \\
 \text{s.t.} \quad & g_1(x) = 2x_1^2 + x_2^2 + x_3^2 + 2x_1 + x_2 + x_4 - 5 \leq 0, \\
 & g_2(x) = x_1^2 + x_2^2 + x_3^2 + x_4^2 + x_1 - x_2 + x_3 - x_4 - 8 \leq 0, \\
 & g_3(x) = x_1^2 + 2x_2^2 + x_3^2 + 2x_4^2 - x_1 - x_4 - 10 \leq 0.
 \end{aligned}$$

Let  $x^0 = (0, 0, 0, 0)$ ,  $\rho_0 = 10$ ,  $N = 4$ ,  $\epsilon_0 = 0.02$ ,  $\eta = 0.01$  and  $\epsilon = 10^{-6}$ . Numerical results of Algorithm I for solving (P3.1) are given in Table 1.

From Table 1, it is said that an approximate  $\epsilon$ -feasible solution to (P3.1) is obtained at the 3'th iteration, that is

$$x^3 = (0.168232, 0.834156, 2.010050, -0.963345)$$

with corresponding objective function value  $f(x^3) = -44.233662$ . It is easy to check that the  $x^3$  is feasible solution to (P3.1). The solution we obtained is slightly better than the solution obtained in the 4'th iteration by method in [8] (the objective function value  $f(x^*) = -44.23040$ ) for this example. Further, with the same parameters  $\rho_0$ ,  $N$ ,  $\epsilon_0$ ,  $\eta$  as above, we change the starting point to  $x^0 = (1, 1, 1, 1)$  or  $x^0 = (6, 6, 6, 6)$ . New numerical results by Algorithm I are given in Table 2 and Table 3.

It is easy to see from Tables 2 and 3 that the convergence of Algorithm I is the same and the objective function values are almost the same. That is to say, the efficiency of Algorithm I does not completely depend on the starting point  $x^0$ . Then, we can choose any starting point for Algorithm I.

*Note:*  $j$  is the number of iteration in the Algorithm I.

$\rho_j$  is constrain penalty parameter at the  $j$ 'th iteration.

$x^j$  is a solution at the  $j$ 'th iteration in the Algorithm I.

$f(x^j)$  is an objective value at  $x^j$ .

$g_i(x^j)$  ( $i = 1, \dots, m$ ) is a constrain value at  $x^j$ .

Table 1: Numerical results of Algorithm I with  $x^0 = (0, 0, 0, 0)$ ,  $\rho_0 = 10$ ,  $N = 4$ 

j	$\rho_j$	$\epsilon_j$	$f(x^j)$	$g_1(x^j)$	$g_2(x^j)$	$g_3(x^j)$	$x^j$
1	10	0.02	-44.271512	0.009467	0.001542	-1.866763	(0.170216, 0.836027, 2.010594, -0.966369)
2	40	0.0002	-44.234025	0.000047	0.000077	-1.883044	(0.169563, 0.835533, 2.008644, -0.964884)
3	160	0.000002	-44.233662	-0.000000	-0.000072	-1.888579	(0.168232, 0.834156, 2.010050, -0.963345)

Table 2: Numerical results of Algorithm I with  $x^0 = (1, 1, 1, 1)$ ,  $\rho_0 = 10$ ,  $N = 4$ 

j	$\rho_j$	$\epsilon_j$	$f(x^j)$	$g_1(x^j)$	$g_2(x^j)$	$g_3(x^j)$	$x^j$
1	10	0.02	-44.271512	0.009467	0.001542	-1.866763	(0.170216, 0.836027, 2.010594, -0.966369)
2	40	0.0002	-44.234025	0.000047	0.000077	-1.883044	(0.169563, 0.835533, 2.008644, -0.964884)
3	160	0.000002	-44.233355	-0.000113	-0.000079	-1.900244	(0.166329, 0.831255, 2.012529, -0.960615)

**Example 2.** Consider the example in [18],

$$\begin{aligned}
 \text{(P3.2)} \quad & \min \quad f(x) = -2x_1 - 6x_2 + x_1^2 - 2x_1x_2 + 2x_2^2 \\
 & \text{s.t.} \quad x_1 + x_2 \leq 2, \\
 & \quad \quad -x_1 + 2x_2 \leq 2, \\
 & \quad \quad x_1, x_2 \geq 0.
 \end{aligned}$$

Let  $g_1(x) = x_1 + x_2 - 2$ ,  $g_2(x) = -x_1 + 2x_2 - 2$ ,  $g_3(x) = -x_1$ ,  $g_4(x) = -x_2$ . Thus problem (P3.2) is equivalent to the following problem:

$$\begin{aligned}
 \text{(P3.2')} \quad & \min \quad f(x) = -2x_1 - 6x_2 + x_1^2 - 2x_1x_2 + 2x_2^2 \\
 & \text{s.t.} \quad g_1(x) = x_1 + x_2 - 2 \leq 0, \\
 & \quad \quad g_2(x) = -x_1 + 2x_2 - 2 \leq 0, \\
 & \quad \quad g_3(x) = -x_1 \leq 0, \\
 & \quad \quad g_4(x) = -x_2 \leq 0.
 \end{aligned}$$

Table 3: Numerical results of Algorithm I with  $x^0 = (6, 6, 6, 6)$ ,  $\rho_0 = 10$ ,  $N = 4$ 

j	$\rho_j$	$\epsilon_j$	$f(x^j)$	$g_1(x^j)$	$g_2(x^j)$	$g_3(x^j)$	$x^j$
1	10	0.02	-44.271512	0.009467	0.001542	-1.866763	(0.170216,0.836027, 2.010594,-0.966369)
2	40	0.0002	-44.234025	0.000047	0.000077	-1.883044	(0.169563,0.835533, 2.008644,-0.964884)
3	160	0.000002	-44.232449	-0.000613	-0.000188	-1.856917	(0.159767,0.840231, 2.011450,-0.963346)

Let  $x^0 = (0, 0)$ ,  $\rho_0 = 8$ ,  $N = 10$ ,  $\epsilon_0 = 0.01$ ,  $\eta = 0.01$  and  $\epsilon = 10^{-6}$ . Numerical results of Algorithm I for solving (P3.2') are given in Table 4.

By Table 4, an approximate optimal solution to (P3.2') is obtained at the 3'th iteration, that is  $x^* = (0.800000, 1.200000)$  with corresponding objective function value  $f(x^*) = -7.200000$ . The solution we obtained is similar with the solution obtained in the 4'th iteration by method in [18] (the objective function value  $f(x^*) = -7.2000$ ) for this example.

Table 4: Numerical results of Algorithm I with  $x^0 = (0, 0)$ ,  $\rho_0 = 8$ ,  $N = 10$ 

j	$\rho_j$	$\epsilon_j$	$f(x^j)$	$g_1(x^j)$	$g_2(x^j)$	$x^j$
1	8	0.01	-7.228666	0.010245	-0.397951	(0.806147,1.204098)
2	80	0.0001	-7.200091	0.000032	-0.399994	(0.800019,1.200013)
3	800	0.000001	-7.200000	0.000000	-0.400000	(0.800000,1.200000)

### 3 Conclusions

This paper has presented a second-order smoothing approximation to the  $l_1$  exact penalty function and an algorithm based on this smoothed penalty problem. It is shown that the optimal solution to the  $(SP_{\rho,\epsilon})$  is an approximate optimal solution to the original optimization problem under some mild conditions. Numerical results show that the Algorithm I has a better convergence for a global approximate solution.

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